# On a way to visualize (some) sheaves 

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#### Abstract

This is an attempt to connect the way in which sheaves are presented in [LM92] with the approach "for children" from [PH1], [PH2], and [FavC] — but these notes are a work in progress that is still in a very preliminary form.


## 1 Sheaves

The archetypal example of a sheaf is the operation

$$
\begin{aligned}
\mathcal{C}^{\infty}: \mathcal{O}(\mathbb{R})^{\mathrm{op}} & \rightarrow \text { Set } \\
U & \mapsto\left\{f: U \rightarrow \mathbb{R} \mid f \text { is } \mathcal{C}^{\infty}\right\}
\end{aligned}
$$

that expects open sets of $\mathbb{R}$ and returns sets of functions; more precisely, for each subset $U \subseteq \mathbb{R}$ it returns $\mathcal{C}^{\infty}(U)$, the set of infinitely differentiable functions from $U$ to $\mathbb{R}$.

This $\mathcal{C}^{\infty}$ is a contravariant functor. The topology $\mathcal{O}(\mathbb{R})$ is a preorder category whose morphisms are the inclusions, and the image by $\mathcal{C}^{\infty}$ of each
inclusion map $V \hookrightarrow U$ is a restriction map $\mathcal{C}^{\infty}(U) \rightarrow \mathcal{C}^{\infty}(V)$. In a diagram:

$\mathcal{O}(\mathbb{R})$

$$
\mathcal{O}(\mathbb{R})^{\mathrm{op}} \xrightarrow{\mathcal{C}^{\infty}} \text { Set }
$$

This diagram follows the conventions from [FavC, section 2]. In brief:
(CAI) the diagram $\emptyset \hookrightarrow V \hookrightarrow U \hookrightarrow \mathbb{R}$ is drawn above $\mathcal{O}(\mathbb{R})$ to indicate that it is inside the category $\mathcal{O}(\mathbb{R})$,
(CFSh) the image of $\emptyset \hookrightarrow V \hookrightarrow U \hookrightarrow \mathbb{R}$ by the functor $\mathcal{C}^{\infty}$ is drawn as a diagram with the same shape as the original; in particular, the unnamed arrow $\mathcal{C}^{\infty}(V) \leftarrow \mathcal{C}^{\infty}(U)$ is the image by $\mathcal{C}^{\infty}$ of the (unnamed) inclusion map $V \hookrightarrow U$,
$(\mathrm{C} \mapsto)$ the arrow $\left.f_{U} \mapsto f_{U}\right|_{V}$ is the internal view of the unnamed arrow $\mathcal{C}^{\infty}(U) \rightarrow \mathcal{C}^{\infty}(V)$, and $\left.f_{U}\right|_{V}:=\mathcal{C}^{\infty}(U \hookleftarrow V)\left(f_{U}\right)$.

The rationale for having an ' $\mathcal{O}(\mathbb{R})^{\prime}$ ' above the ' $\mathcal{O}(\mathbb{R})^{\text {op }}$ ' is explained in [FavC, section 7.4], and our reasons for drawing topological spaces with the "everything" on top and the "nothing" at the bottom are explained in [PH1]; in short, that's because we will at some point treat $\mathcal{O}(\mathbb{R})$ as a logic, in which $\mathbb{R}$ is 'true', 'top', and ' $T$ '.

### 1.1 The unique glueing property

$\mathcal{C}^{\infty}$ has the "unique glueing property". The UGP can be formalized in several different, and slightly incompatible, ways.

Take any two open sets $U_{1}, U_{2} \in \mathcal{O}(\mathbb{R})$ and choose functions $f_{U_{1}} \in \mathcal{C}^{\infty}\left(U_{1}\right)$ and $f_{U_{2}} \in \mathcal{C}^{\infty}\left(U_{2}\right)$. We say that $f_{U_{1}}$ and $f_{U_{2}}$ are compatible when their
restrictions to $\mathcal{C}^{\infty}\left(U_{1} \cap U_{2}\right)$ are the same - i.e., when $\left.f_{U_{1}}\right|_{U_{1} \cap U_{2}}=\left.f_{U_{2}}\right|_{U_{1} \cap U_{2}}$. In a diagram:


Our first version of the unique glueing property is this. For all $U_{1}, U_{2} \in$ $\mathcal{O}(\mathbb{R})$, for every compatible pair $\left(f_{U_{1}}, f_{U_{2}}\right) \in \mathcal{O}\left(U_{1}\right) \times \mathcal{O}\left(U_{2}\right)$ has a unique glueing: there is a unique $f \in \mathcal{C}^{\infty}\left(U_{1} \cup U_{2}\right)$ such that this $f$ restricts to $f_{U_{1}}$ and $f_{U_{2}}$, i.e., $\left.f\right|_{U_{1}}=f_{U_{1}}$ and $\left.f\right|_{U_{2}}=f_{U_{2}}$. In a diagram:


Our second version of the UGP deals with compatible families of functions. Take an index set $I$, a family of open sets $\left(U_{i}\right)_{i \in I}$ such that each $U_{i} \in \mathcal{O}(\mathbb{R})$, and a family of functions $\left(f_{i}\right)_{i \in I}$ such that each $f_{i} \in U_{i}$. We say that this family of functions is pairwise compatible if

$$
\forall i, j \in I .\left.f_{i}\right|_{U_{i} \cap U_{j}}=\left.f_{j}\right|_{U_{i} \cap U_{j}},
$$

and we say that this family has a unique glueing if there is a unique $f \in$ $\mathcal{C}^{\infty}(U)$, where $U:=\bigcup_{i \in I} U_{i}$, such that

$$
\forall i \in I . f_{i}=\left.f\right|_{U_{i}}
$$

A standard reference for this is [LM92], section II.1, but it uses some notational tricks that took me years (really!!!) to understand... they are easy to decypher when we know some notations from dependent types, though.

Let $A=\{2,3\}, B=\{4,5\}, C_{2}=\{20,21\}, C_{3}=\{32,33\}$. The notations for the set of all pairs $(a, b)$ with $a \in A$ and $b \in B$ and for all function that take each $a \in A$ to a $b \in B$ are well-known: $A \times B$ and $B^{A}$, but we will sometimes write $B^{A}$ as $A \rightarrow B$. The standard notations for the set of
all pairs ( $a, c$ ) with $a \in A$ and $c \in C_{a}$ and for the set of all functions that take each $a \in A$ to a $c \in C_{a}$ are less familiar (see [HOTT, sections 1.4 and 1.6]): $\Sigma_{a \in A} \cdot C_{a}$ and $\Pi_{a \in A} \cdot C_{a}$. Some programming languages with support for dependent types, such as Agda, implement a notation that looks like an extension of ' $x$ ' and ' $\rightarrow$ ':

$$
\begin{aligned}
(a: A) \times C_{a} & :=\Sigma_{a \in A} \cdot C_{a} \\
(a: A) \rightarrow C_{a} & :=\Pi_{a \in A} \cdot C_{a}
\end{aligned}
$$

So we have:

$$
\begin{aligned}
(a: A) \times C_{a}:= & \{(2,20),(2,21),(3,32),(3,33)\}, \\
(a: A) \rightarrow C_{a}:= & \{\{(2,20),(3,32)\}, \\
& \{(2,20),(3,33)\}, \\
& \{(2,21),(3,32)\}, \\
& \{(2,21),(3,33)\}\} .
\end{aligned}
$$

A family of pairwise functions on $\mathcal{C}^{\infty}$ is a triple

$$
(I, \mathcal{U}, \mathcal{F}) \quad: \quad(I: \text { Sets }) \times(\mathcal{U}: I \rightarrow \mathcal{O}(\mathbb{R})) \times\left(\mathcal{F}:(i: I) \rightarrow \mathcal{C}^{\infty}(\mathcal{U} i)\right)
$$

obeying:

$$
\forall i, j \in I .\left.(\mathcal{F} i)\right|_{(\mathcal{U i} \cap \mathcal{U} j)}=\left.(\mathcal{F} j)\right|_{\left(\mathcal{U} i \cap \mathcal{U}_{j}\right)} .
$$

Our third version of the UGP needs "downward-closedness". For a subset $\mathcal{V} \subseteq \mathcal{O}(\mathbb{R})$ we define

$$
\downarrow \mathcal{V}:=\{W \in \mathcal{O}(\mathbb{R}) \mid \exists V \in \mathcal{V} . W \subseteq V\}
$$

and we say that a $\mathcal{V} \subseteq \mathcal{O}(\mathbb{R})$ is sieve when $\mathcal{V}=\downarrow \mathcal{V}$. A compatible family $\mathcal{G}$ on a sieve $\mathcal{V}$ is a family $\mathcal{G}:(V: \mathcal{V}) \rightarrow \mathcal{C}^{\infty}(V)$ obeying:

$$
\forall V, W \in \mathcal{V} . W \subseteq V \rightarrow \mathcal{G} W=\left.\mathcal{G} V\right|_{W}
$$

A compatible family $\mathcal{G}$ on a sieve $\mathcal{V}$ has a unique glueing iff there is a unique $f \in \mathcal{C}^{\infty}(U)$, where $U=\bigcup \mathcal{V}$, such that

$$
\forall V \in \mathcal{V},\left.f\right|_{V}=\mathcal{G} V
$$

A good way to understand how these ideas can be generalized is to work on cases where everything is finite and everything can be drawn explicitly.

### 1.2 Presheaves (on some finite topologies)

In the previous sections we worked with a fixed topology, $\mathcal{O}(\mathbb{R})$, and with a fixed contravariant functor, $\mathcal{C}^{\infty}: \mathcal{O}(\mathbb{R})^{\text {op }} \rightarrow$ Set. We saw that $\mathcal{C}^{\infty}$ obeys three different "unique glueing properties" that were easy to understand; we will now generalize this, in two steps.

Let $X$ be the open interval $(0,3) \subset \mathbb{R}$, and let

$$
\mathcal{O}(X):=\{\emptyset,(1,2),(0,2),(1,3),(0,3)\},
$$

which is a subtopology of the usual topology on $(0,3)$. As $\mathcal{O}(X)$ is finite we can draw it, and its image by $\mathcal{C}^{\infty}$, explicitly:


In [PH1, sections 1, 2, 12, and 13], we saw how to interpret diagrams like -• as directed acyclical graphs (DAGs), how to define an order topology on a DAG, and how to draw these topologies. If we replace $X$ by $\bullet \bullet$ and $\mathcal{C}^{\infty}$
by an arbitrary contravariant functor $F: \mathcal{O}(\bullet \bullet)^{\mathrm{op}} \rightarrow$ Set we get this,

that by (an intentional) coincidence has the same shape as the previous diagram. The trick to "pronounce" things like ${ }^{0} 1^{1}$ is explained in [PH1, section 1]: if we read aloud the digits of ${ }^{0}{ }^{1}$ in "reading order", i.e., for top to bottom and in each line from left to right, then it becomes "zero-one-one".

Now that our topology has a definite shape we can use that shape, with ' 0 's and ' 1 's at the right positions, to talk of subsets of it. For example, $\left\{1_{1}{ }^{0},{ }^{0} 1^{1}\right\}={ }_{1}^{0}{ }_{0}^{0}$, and this is not a sieve, because $\downarrow{ }_{0}^{1}{ }_{0}^{0}{ }_{0}^{1}={ }_{1}^{1}{ }_{1}^{0}{ }_{1}^{1}$; also, $\bigcup{ }_{0}^{0}{ }_{0}^{1}=$ $\left.\bigcup\left\{1_{1}{ }^{0},{ }_{0}{ }_{1}\right\}\right\}={ }_{1}{ }_{1}{ }^{1}$, and $\downarrow\left\{\bigcup{ }_{1}^{1}{ }_{0}^{0}{ }_{0}^{1}\right\}={ }_{0}^{1}{ }_{1}^{1}{ }_{1}$.

We can also choose other presheaves and test if they obey the unique glueing properties. Let $E$ be this functor:


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For the first version of the UGP, let $U_{1}={ }^{1}{ }_{1}{ }^{0}, U_{2}={ }^{0}{ }_{1}{ }^{1}, f_{U_{1}}=1, f_{U_{2}}=2$. This is a compatible pair, because $U_{1} \cap U_{2}={ }^{0}{ }^{0}$ and $\left.f_{U_{1}}\right|_{U_{1} \cap U_{2}}=\left.f_{U_{2}}\right|_{U_{1} \cap U_{2}}=1$ - but there is no $f \in E\left(U_{1} \cup U_{2}\right)$ that restricts to both $f_{U_{1}}$ and $f_{U_{2}}$; to check this we just need to test the two candidates, 23 and 24.

The second version of the UGP needs an index set. Take $I=\{42,200\}$ 42 and 200 are two of my favorite numbers -, and $\mathcal{U}=\left\{\left(42,{ }_{1}{ }^{0}\right),\left(200,{ }^{0} 1^{1}\right)\right\}$, $\mathcal{F}=\{(42,1),(200,2)\}$. This is a pairwise compatible family of functions, but it has two possible glueings. We have $U:=\bigcup_{i \in I} U_{i}={ }^{1}{ }_{1}{ }^{1}$, and both 23 and 24 are glueings, i.e., are possible values for the $f \in E\left(\bigcup_{i \in I} U_{i}\right)$ that obeys $\forall i \in I . f_{i}=\left.f\right|_{U_{i}}$.

The third version of the UGP is the nicest to draw. A pair $(\mathcal{V} \subseteq$ $\left.\mathcal{O}\left({ }^{1} 1^{1}\right), \mathcal{G}:(V: \mathcal{V}) \rightarrow E(V)\right)$ can be drawn using the conventions for drawing partial functions from [PH1, section 1]: we call $\mathcal{V}$ the support of the family $(\mathcal{V}, \mathcal{G})$, and for each $V \in \mathcal{O}(\bullet \cdot \bullet)$ we draw $\mathcal{G} V$ on the position of $V$ if $V \in \mathcal{V}$; when $V \notin \mathcal{V}$ we draw a $‘$ '. For example, let

$$
\begin{aligned}
& \mathcal{V}=\left\{{ }_{1}{ }_{1}{ }^{0},{ }_{1} 1_{1}{ }^{1},{ }_{0}{ }_{0}{ }^{0}\right\}=\bullet \bullet=\begin{array}{c}
1_{0}^{0}{ }_{1}{ }_{1}, ~ \\
1
\end{array} \\
& \mathcal{G}=\left\{\left({ }_{1}{ }^{0}, 1\right),\left({ }^{0}{ }_{1}{ }^{1}, 2\right),\left({ }_{0}{ }_{0}{ }^{0}, 0\right)\right\}={ }^{1} \dot{0}^{2}, \\
& (\mathcal{V}, \mathcal{G})=\left(\because, \quad \begin{array}{cc}
\dot{0} & 2 \\
\bullet & \\
0
\end{array}\right.
\end{aligned}
$$

This $\mathcal{V}=\bullet$ isn't a sieve, and this $\mathcal{G}={ }^{1} \dot{0}^{2}$ isn't compatible downwards. We want to restrict our attention to sieves, and it turns out the set of all sieves on $\mathcal{O}(\bullet \bullet \bullet)$ is exactly $\mathcal{O}(\mathcal{O}(\bullet \bullet))$, the order topology on the partial order $\because \bullet$ A very good place to learn about this is the section about down-sets and up-sets of [DP02, page 20 onwards]; for an ordered set $P$ they define $\mathcal{O}(P)$ as the "ordered set of down-sets of $P$ " in a way that is very easy to iterate

$$
\begin{aligned}
& \text { ヶ } \downarrow \downarrow \\
& \left(\begin{array}{c}
0 \\
1_{1} \\
1 \\
1
\end{array}\right) \mapsto\left\{\begin{array}{cc}
1 & 1 \\
1 & 1 \\
1 & 1 \\
1 & 1 \\
1 & 2
\end{array}\right\} \quad\left(\begin{array}{c}
0 \\
0 \\
0 \\
0 \\
0
\end{array}\right) \longmapsto\{\quad \vdots\} \\
& \operatorname{Sub}\left({ }^{1}{ }_{1}^{1} 1\right. \\
& \operatorname{Sub}\left({ }^{1}{ }_{1}^{1} 1\right) \\
& \left.\operatorname{Sub}\left({ }_{(1}^{1}{ }_{1}^{1}\right)^{1}\right)^{\text {op }} \xrightarrow{E} \operatorname{Set} \quad \operatorname{Sub}\left({ }_{1}^{1}{ }_{1}^{1}\right)_{1}^{\mathrm{op}} \xrightarrow{E} \operatorname{Set}
\end{aligned}
$$



## References

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[PH1] E. Ochs. "Planar Heyting Algebras for Children". In: South American Journal of Logic 5.1 (2019). http://angg.twu.net/math-b.html\#zhas-for-children-2, pp. 125-164.
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